

Figure 1: Transaction Processing System (Prior Art).

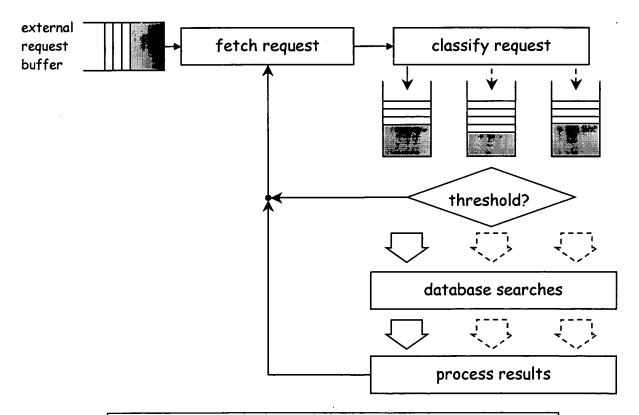


Figure 2: Transaction Processing System with Request Buffering.



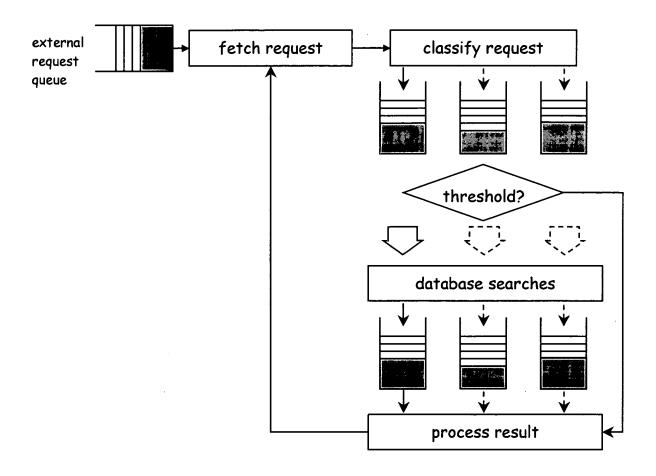
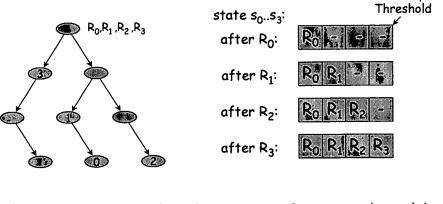


Figure 3: Transaction Processing System with Request and Result Buffering.



## First Set of Search Requests



node location of key for R<sub>i</sub>
 (not the key value itself)

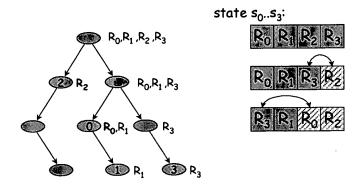
 $\bigcirc$  R<sub>j</sub>,R<sub>k</sub>: current node of R<sub>j</sub>,R<sub>k</sub>

R<sub>i</sub>: requested search key, current node pointer, tag (or thread id), etc.

Startup

Figure 4: Example of a tree traversal buffering.

## First Pipelined Search



prefetch s<sub>0</sub>..s<sub>3</sub>
while pending > min
loop i from 0 to3:
work R<sub>i</sub>
update state s<sub>i</sub>
prefetch s<sub>i</sub>

Figure 5: Example of a pipelined tree search traversal.



## Second Pipelined Search

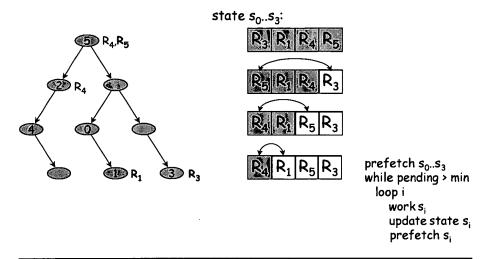
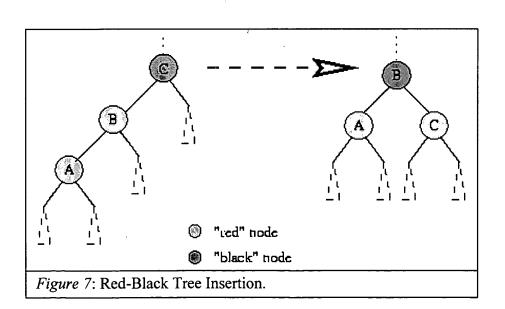


Figure 6: Example of a pipelined tree search traversal state.





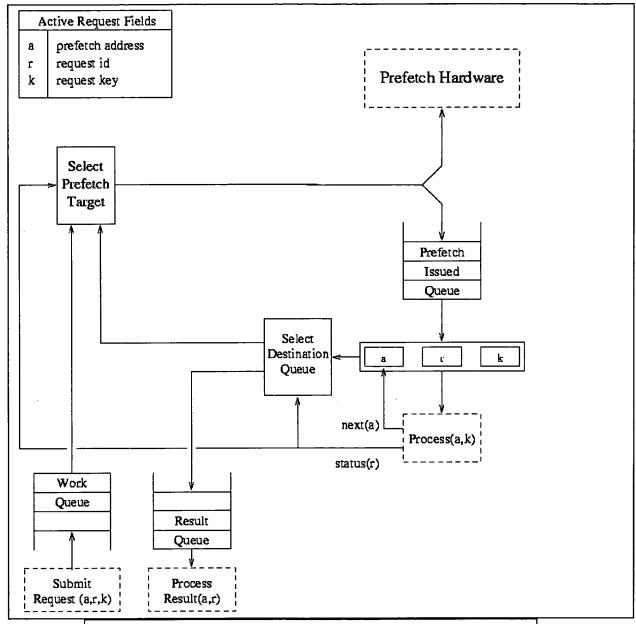


Figure 8: Restructuring mechanism, as implemented in software.



```
RESTRUCTURED-TRAVERSAL( S, request )

begin

AQ.enqueue( request );

if AQ.size ≥ K then

SOFTWARE-PIPELINE( S, AQ, RQ );

if RQ.size = 0 then

return POSTPONE

else

return RQ.dequeue()

end
```

Figure 9: Accumulating K requests on accumulation queue AQ for software pipelined traversals of data structure S, where K is the startup threshold. Accumulated results are turned from result queue RQ.

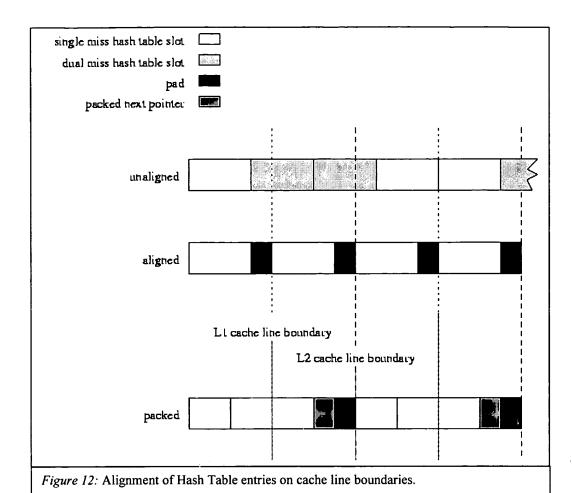


```
TREE-DELAYED-SEARCH(lower)
begin
         integer i, prologue;
        prologue \leftarrow MIN(lower, RQ.size);
         i \leftarrow 0;
         while i<prologue do
                  PREFETCH( RQ.elem[i]);
                  i \leftarrow i + 1;
         end while
         TREE-RECURSIVE-SEARCH( lower );
end
Figure 10: Recursive search requests, initial pre-recursive component.
```

```
TREE-RECURSIVE-SEARCH(lower)
begin
         i \leftarrow 0;
         while i<AQ.size do
                   request \leftarrow AQ.elem[i];
                   k \leftarrow request.key;
                   n \leftarrow request.node;
                   if n = NIL or k = n.key then
                             AQ.delete( request );
                             RQ.enqueue(request);
                   else
                             if k < n.key then request.node \leftarrow n.left;
                                           else request.node \leftarrow n.right;
                             endif
                             PREFETCH( request.node );
                   endif
                   i \leftarrow i + 1;
          if AQ.size \ge lower then TREE-RECURSIVE-SEARCH( lower ); endif
end
```

Figure 11: Recursive search requests, recursive component.







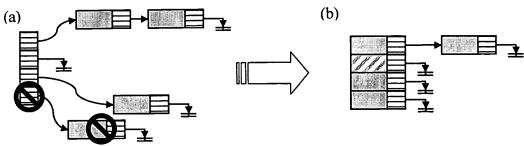
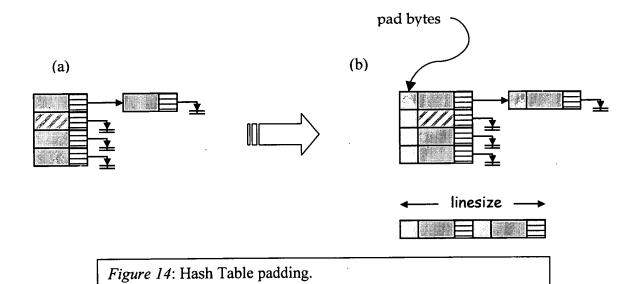


Figure 13: Hash Table homogenezation.



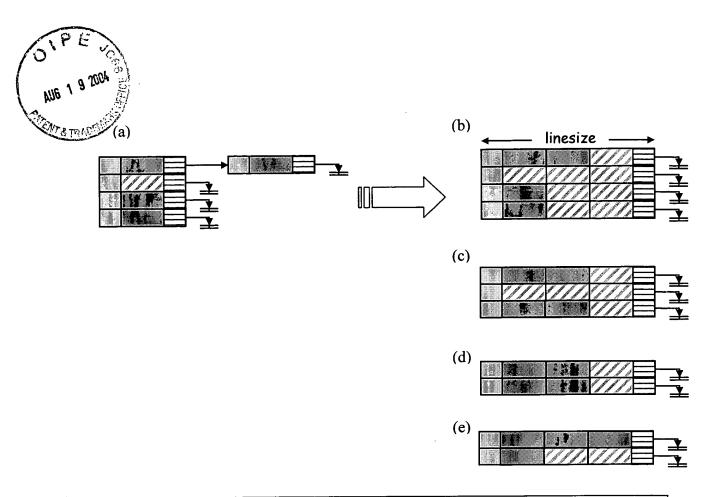


Figure 15: Hash table packing. Representing a homogeneous hash table structure (a) as a packed structure (b), which can be re-balanced to make the table less sparse as in (c), (d), or (e).



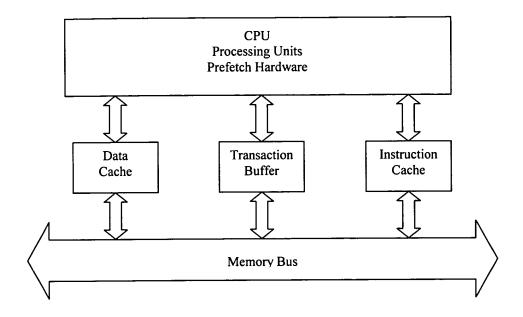


Figure 16: Transaction Buffer.



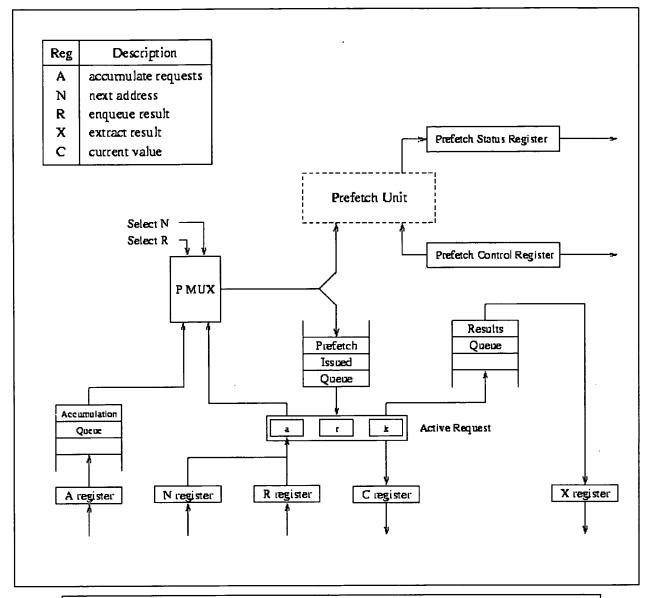


Figure 17: Transaction Buffer Details, single set of queues.